AComponentFrameworkSupportingPeerServicesfor SpaceDataManagement

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toincreasethecollaborationandcorrelationofkey

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Abstract—Knowledgediscoveryanddatacorrelation requireaunifiedapproachtobasicdatamanagement. However, achieving such an approach is nearly impossible withhundredsofdisparatedatasources, legacysystems and dataformats. This problem is pervasive inthespacescience communitywheredatamodels,taxonomiesanddata managementsystemsarelocallyimplementedandlimited metadatahasbeencollectedandorganized.Technology developedbytheObjectOrientedDataTechnology (OODT)taskatNASA'sJet PropulsionLaboratory(JPL) hasbeen exploring component frameworks formanging, locatingandexchangingdataresidingwithina geographically distributed network. OOD Thas taken a novelapproachtowardssolvingthisproblembyexploiting webtechnologie susuallydedicatedtoe combinedwitharich.metadata -basedenvironment.The componentsdevelopedbyOODTcreateasetofdistributed peer-to-peerservicesthatallowfordatamanagedbyapeer tobesearchedandreturnedaspartofanintegrat eddata managementsystem. This paper discusses the approach takentodevelopasoftwareframework, and two prototype development efforts for the Planetary Data System (PDS)and the Mission and Ground Asset Database.

TABLEOF CONTENTS

- 1. Introduction
- 2. HISTORYAND PROGRAMMATIC GOALS
- 3. PEER SERVICES ARCHITECTURE
- 4. Information Architecture
- 5. APPLICATIONTOTHE PLANETARY DATA SYSTEM
- 6. APPLICATIONTOTHE DEEP SPACE NETWORK
- 7. Future
- 8. References
- 9. ABOUTTHE AUTHORS

1. Introduction

TheObjectOrientedDataTechnology(OODT)task, supportedbytheNationalAeronauticsandSpace Administration(NASA)andimplementedattheJet PropulsionLaboratory(JPL),isfocusingondevelo pinga referenceframeworktosupportcrossdisciplinarysolutions fordataaccess,exchangeanddistributionofdistributeddata

engineeringandsciencedataproducedthroughoutthe lifecycleofaspace mission.NASAisambitiouslyplanning toflyseveralnewmissionswithanincreasingamountof dataplannedforeachmission. Unfortunately, the data collectedfromcradletograveofthesemissionstendstobe difficulttolocate, access, understandand interoperability,softwareanddatareuse,knowledge discoveryandincreasingmissioncosts. Asaresult, OODT identifiedseveralkeygoalsincluding(1)understanding wherethedataresourcesreside,(2)understandinghowdata isaccessedi neachsystem,(3)understandingtheunderlying datamodelsforeachsystem,and(4)buildingubiquitous interfacesacrossmultipledatasystemsthatdemonstrate interoperabilityusingacommonquerymechanism.In ordertomeetthesechallenginggoalswe definedan architectureandasetofrequirements. Thearchitecture definedaframework, services, and are ference information modelthatwastestedagainstseveralprojectsincludingthe PlanetaryDataSystem(PDS)andtheDeepSpaceNetwork (DSN). We based requirements on experience from these activities as well as interactions with several other efforts whichsupporttheoverallobjectivesandcultureofJPLand NASA. The partnerships established by the OOD Tteam to researchsolutionsprovedtobecr iticaltomeetingourgoals, andallowedustoidentifykeydataarchitecturesthat supportthemanagementandexchangeofdatainaneraof ever-increasing disparated at a sources.

2. HISTORYAND PROGRAMMATIC GOALS

InFebruary1998,JPLwasfundedbytheOffi ScienceatNASAtoperformresearchinobjectorientedand distributeddatamanagementtechnologies.Aprincipalgoal oftheOODTtaskistoexploreanddevelopcutting -edge systemsfortheinteroperabilityofscientificresearch, measurements, and data that span multiple disciplines. Giventheamountofdatageneratedbypastunmanneddeep spacemissions, and the predicted quantity of data that is plannedforfuturemissions, keyinformation technology solutionsneedtobecreatedthatwille anddistributionofsciencedataproductsresidinginmultiple datasystems. As a result, the focus of OOD Thas been to provideanarchitecturalframeworkthatincludessoftware servicesfordataarchiving,datadistribution,dataloca tion anddataexchange. The ability to correlate data across multiplemissions, databases and datasets is dependent on havingthedatamanagementinfrastructuretoallowthe

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exchangeofdata. Advances intechnologies including the ExtensibleMarkupLang uage(XML) [1],metadata management, and distributed computing have allowed such aninfrastructuretobeareality.

Presently, spacescientists cannot easily locateor used at a acrossthehundredsofautonomous,heterogeneousa nd distributeddatasystemswithinthespacescience community. The various platforms, data formats, data interpretations and Database Management Systems (DBMS) usedmakeinteroperabilityadifficulttask.Developinga ubiquitouscomputingenvironmentac rossdatasystemsisa difficultchallengegiventheindependenceofworkgroups andresearchactivitiesatNASA. Therefore, the architecture thatwehavedefinedallowseachdatasystemtobe independentlydeveloped,however,thesystemsexchange datagi venamessaginginfrastructurethatincludes metadata²descriptions associated with the messages. The architectureisbasedonadirectedgraph(ordigraph)ofdata systems that are traversed using a common query messageimplementedinXMLinordertosat isfyaquery.Profiles setsofresourcedefinitions -describenodes(ordata systems)thatarepartofthedigraph.Profilesmaypointto otherprofilesthusrepresentingarcsofthedigraph.A profileisessentiallyametadatadescriptionoftheresourc knownatanodeinthedistributedframework. These resourcescanbedefinitionsofdatasystems, ordata productsmanagedbyadatasystem.Boththeserviceand informationarchitecturewillbedefinedinmoredetaillater inthepaper.

Assystemsbe comemoredistributed, theneed to integrate dissimilar systems increases. A simple calculation yields thatanorganizationwith **N**numberofsystemscould $(N^2-N)/2$ interfacesin potentiallyidentifytheneedtocreate ordertosharedatabetweenallofth osesystems.Asa result, the need to identify and supports tandards based interfacesincreasesasthenumberofdistributednodesinan organizationincreases. One of the emerging solutions is peer-to-peer(P2P)architectures.AP2Parchitectureis definedas"environmentconsist[ing]ofcomputerswith equalcapabilities that shareresources (such as processing powerandmemory), communicate exclusively with each otheranddonotconnecttoserversorcentraldatabases" OneofthekeygoalsofOODThasbeentoprovidethe abilityfordatasystemstobecomepeersthatprovide services. These services could distribute information such assciencedata, engineering data, navigation data or even containcomputationalcapabi litiessuchasscience processing algorithms or simulation environments. The Internethasspawnedseveralsuchcapabilitiesincluding Napsterforsharingmusicfiles. It is important to note that whileitisgenerallyacceptedthatfederatingtogetherda systemswillincreasetheknowledgegeneratedby correlatingdatacapturedindispersedspacescience

²Metadataisdataaboutdata.

databases, it is still unclear as to what the return on investmentwillyield.SeveralinitiativesatNASAare exploringthatinadditiontoOODTincl udingtheNational VirtualObservatory(NVO)whichisattemptingtobuilda federatedenvironmentforcollaborationandresearchthat integratesastrophysicsandastronomydatabasesacrossthe sciencecommunitywithagoalto"enablenewscience, moreeffe ctivescience, and more cost effective science" [3]. Asaresult, one of the activities outlined by NVO is to "[establish]acommonsystemsapproachtodatapipelining, archiving, and retrieval that will ensure easy access by a largeanddiversecommunityofusersandthatwillminimize costsandtimestocompletion." [3]Astheseefforts continueonegoalisclear:theneedtohaveageneric, adaptable, scalable and extensible architecture with wel 1 definedinterfacesthatimplementstandardsbasedprotocols andarichmetadatainfrastructureisanimportantkeyto unlockinginteroperabilityacrossNASA.

3. PEER SERVICES ARCHITECTURE

OODTconsistsofasetofcooperating, distributed peer components. Althoughitscurrentincarnationis implemented with a client -server communications substrate, thedesignresemblesaP2Pnetwork,andwehaveplansto transitionthecommunicationssubstratetoaP2P implementationshortly [4].

DistributedFramework

es

[2].

ta

OODTisadistributedsystem, wherein components may be geographicallydispersed. Astandard Internet TCP/IP connectionsufficestoprovideconnectivitybetween components at this time. In addition, we've architected the software asaplug -insystem [5].RatherthanmakeOODT asetofutilityAPIsoutofwhicheachsitedevelopsitsown applications, we create a set of classes and interfaces that sitescustomizeandimplementandthenregisterwiththe OODTAPI.Frameworksareeasiertousebecausethe frameworkclassesandinterfacestellthedeveloperexactly whatneedstobeimplemented. Moreworknecessarily falls upontheframeworkdevelopers.

WithinOODTtherearethreemajorcomponents:

- **Profileservers** servescientificmetadataandcantell whetheraparticularresourcecanprovideananswertoa query.
- **Productservers** servedataproductsinasystem independentformatwhenpresentedwithaproduct request.
- Queryservers acceptprofileandprodu ctqueriesand traversethenetworkofprofileandproductservers, collectingresults.

Thequeryserveristhemainentrypointintothesystem, withApplicationProgrammer'sInterface(API)for

developersandawebinterfaceforendusersandtesting.I additiontothesethreeservers, anarchiveserverprovidesa metadata-capabledatastoragesystemforanykindof productthat can be represented as a sequence of bytes (such a safile).

CommonDesignElements

Profileandproductservers, in addition to being distributed peer components, exhibits ome common designelements. First, they all advertise are motely accessible interface, call ablevia CORBA [6]. While clients may contact a profile or products erverdirectly, ween courage them to use the query service which simulates the gateway to the P2P space. Figure 1, a UML [7] deployment diagram, demonstrates a sample deployment of the OODTs of tware. In that diagram, the query server has network connections to two profile and two products ervers.

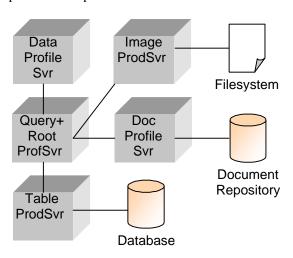


Figure 1 -SampleOODTDeployment

Also,profileandproductservershaverun -timespecified, replaceablebackendprocessors. Aprofileservermayget itsmetadatafromaflatfileoradatabasebychangingthe backend.Similarly,aproductservermayretrieveanimage fromthefilesystemorfromanimagestoragedevicewitha customizedbackend.Backendinterfacesmandatethe callingconventionsforprofileandproductservers.We achievethisrun -timeconfigurationthroughJava's [8] interfaceanddynamicloadingmechanisms;coreproduct andprofileserversreferonlytoabstractbackendinterfac es, whilespecificnamedclassesimplementthoseinterfaces.

Finally,profileandproductserversallmanipulatethesame querystructure. This structure exists as an XMLD ocument Type Definition and as Java class, and iten capsulates the user's queryp lusary results retrieved so far.

ProfileServer

The profiles erver's jobistos erve profiles [9], which are XML-representable metadata descriptions of resources. A resource in this context may be a single data granule, a dataset (a collection of granules), a collection of datasets, a

datadictionary, anything with a Uniform Resource Identifier (URI) [10], such as a URL, or other profiles in other profiles ervers. A profile "profiles" are source b describing data elements it contains, what data dictionary defines those elements, who created it, when, in what language, under what subject headings, and so forth. Section 4 describes the profile information model completely.

Responsibilities—Theprimaryresponsibilityofaprofile serveristoacceptaqueryanddetermineifanyofthe profilesitmanagesprofileresourcesthatcouldsatisfythat query. Eachprofileserveracceptsthequeryanddelegates toabackendimpl ementationtosearchthesetofmanaged profilesandreturnanythatmatch. Figure 2depictsthe classarchitectureinaUMLdiagram.

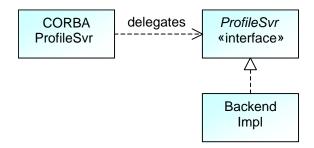


Figure 2 -ProfileServiceClassArchitect ure

Althoughuserscancontactaspecificprofileserverdirectly throughitsCORBAinterface, they typically use the query serviceinstead. Webootstrapthequeryserver with a "root" profileserverwhosesetofmanagedprofilesrefertoother profiles ervers. Thoseotherprofiles ervers may inturn refer toyetmoreprofileservers, forming a directed acyclic graph ofservers. Thequery server can crawlthis graph (eliminating redundant cyclic queries) and gather results fromtheentirenetwork.In Figure 1,thequeryserverruns boththequeryserverandtherootprofileserver;theroot profileservercontainstwoprofilesthatrefertothe documentprofileserverandthedataprofileserver. The scribes URL s in the documentdocumentprofileserverde repository, and the data profiles erver describes the image andtableproductservers.

Besidessearchingasetofprofiles,eachprofileserverhas secondaryresponsibilitiestomanageitssetofprofiles. The managementinter faceincludestypicalsetoperations such as adding, modifying, and deleting profiles.

CurrentBackends — Theprofileserver'splug -inarchitecture hasenabledustodesignandexperimentwithseveral backendimplementationsforqueryingandmanaging profiles. Because profiles can be represented in XML, our original implementation would store each profile in memory as an XML Document Object Model (DOM) tree, and search by traversing the tree.

Another implementation uses an Oracle relational database, essentially transforming the DOM tree into a series of

databasetables. This backend transforms the user's query into SQL and uses a stored procedure to retrieve matching profiles in XML format.

Highlyspecificapplicationsofprofilesarepossible,too.At JPL,Xerox'sDocuShare [11]productmanageslarge collectionsofdocumentation.Aprofileserverbackendthat usesDocuShare'sWebDAV -likeinterface [12]enablesa singlequerytospanmultipleDocuSh areandother documentationmanagementsystems,enhancingknowledge sharing.Inthisbackend,profilesofresourcesarebuilton the-flybyassemblinginformationfromthedocumentation managementsystem.

ProductServer

Productserversexisttoprovidean interfacetoretrieving specificscientificdataproducts. Likeprofileservers, productserversacceptthequerystructurementionedearlier anddescribedinSection 4. Insteadofaddingprofilestothe resultssectionoft hequerystructure, they addactual data products. Dataproducts in this sense may be individual granules, datasets, or collections of datasets, depending on the backen dimplementation to which the products erver delegates its query requests.

Aswithprof ileservers, products ervers delegater equests for products to aspecific backen dimplementation. The backend's jobisto actast hetranslator between a system dependent products to rage system and the system independent OODT framework. It accepts a quer transform sitinto a format required by the product storage system, retrieves the product, transform sitinto a system independent format, and add sitt othe query structure. Figure 3 shows the class architecture.

y,

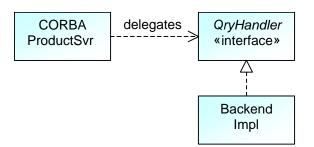


Figure 3 -ProductServiceClassArchitecture

The query handler interface for products ervers is far simpler than the profiles erver interface for profiles ervers.

There is a single method that the backend must implement that accepts the generic query structure and returns it without matching products in stalled.

TypicalOperation —Usersusuallywon'tknowaheadof timewhatproductserversexistsandbywhatnetwork namesthey'reknown.Likewise,thequerys erverdoesnot knowwhatproductserversareavailableeither —it's bootstrappedwithonlyasingleprofileserver,the"root" profileserver.Someprofileserversinthenetworkmust

containmatchingprofilesthatdescribeproductserversthat cananswer arequestfordata.

Asaresult, queryingisatwostepprocess, depicted in the UML interaction diagramin Figure 4. First, the user issues a profile query for the desired data. The query service crawls through the network of profiles ervers and returns any matching profiles. The user then selects a product server from the matching set and re issues the query as a product query directed at a specific server. The query service contacts the product server, which retrieves the product.

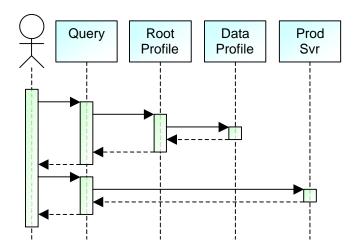


Figure 4 -RetrievingaProduct

ProductFormats —Asmentioned.oneoftheresponsibilities ofthebackendqueryhandlerforaproductserveristo translateproductsfromsystem -dependenttoasyste mindependentformat. These system -independentformatsare anythat can be encoded with an Internet MIME type[13]. Thequerystructure's results section includes the MIME typeoftheproduct.Forexample,aproductservert hat retrievesimagesfromadatabaseinaproprietaryformat maychoosetoreturnthemineithertheimage/jpegor image/pngMIMEformats.Atabularproductmightbein thetext/tab -separated-valuesformat.

TheOODTframeworkprovidesanautomatedmechani sm toencodedtheproductitselfbasedontheMIMEtype.
BecausethequerystructureisserializedintoXMLbetween peercomponents,theframeworkencodescertainMIME typestobecompatiblewithXML'sstringrepresentation.
Forexample,theframeworken codesatext/tab -separated-valuesproductintoanXMLCDATAsection.Itencodes binaryproductslikeimage/giforapplication/ms -wordinto thebase -64textualformat [13].

rs

Usershavesomecontroloverwhatformatproductserve returnproductstothem.Partofthequerystructureincludes alistofacceptableMIMEtypes.Bydefault,thislistis*/*, meaninganyproductformatisacceptable.Bysettingthisto alistofpreferredMIMEtypes,productserverslearnofthe users'desiresandmayencodeproductsappropriately.For example,settingthelisttoimage/jpegandimage/*means

theuserprefersimagesinJPEGformatbutwillacceptany imageformatifnecessary.Settingittoimage/jpegalone meansonlyJPEGformati magesareacceptable.Whena productservercannotsatisfyauser'sformatrequest,itmay returnnoproductevenifitotherwisehadamatchforthe query.

QueryServer

AqueryserverprovidesthepointofentryintotheOODT framework. Thequeryserver contains the algorithms necessary to transform the physical client/server CORBA model into the logical P2P model for executing queries and gathering results. It also provides a simplified interface for users: they don't need to know what profile or produ ct servers exist and what their object addresses are —they only need to access the query server.

Userscanaccessthequeryserverinavarietyofways dependingontheirsophisticationandimplementation language.CORBA -capableuserscanaccessthequery server's CORBA interfacedirectlyandcallmethodsto searchforprofilesorproducts.JavausersmayuseanAPI classthatencapsulatestheCORBAcommunication.Users of other languages may use two different HTTP [14] interfaces, builtusing Java Servlets and servlet -capable web servers [15].

OneHTTPinterfaceiscompletelygeneric:itprovidesthe profileandproductsearchcapabilitiesofthequeryserver throughanHTTPGETorPOSTrequest.T heusercontacts thequeryURL,passingtheXMLquerystructureandother parameters,andreceivingtheXMLquerystructureback, possiblywithencodedresults.Thisinterfacerequiresthat theclientbecapableofprocessingXMLdocuments(and base-64e ncodeddataforbinaryproducts)inorderto retrieveresultsfromthequerystructure.

ThesecondHTTPinterfaceprovidessimplifiedproduct searchesonlyandrequiresnoXMLprocessingcapabilityon theclientside. TheuseragainusesanHTTPGETorPO ST request, passing aquerystring and receives back a single productinits system -independent MIME type.

4. Information Architecture

Spacescienceinstrumentsandexperimentsgeneratescience dataproductsthatarearchivedinsciencedatasystems. Effectivesearchandretrievalofdataproductswithinand acrossthesedatasystemsdependcruciallyoninformation architecturesthatareorganizedonprinciplesknowntothe prospectiveusersanddesignedtoreflectrelevant relationships.Well -designedinform ationarchitecturesalso enabledatabrowsing,datamining,andcorrelativescience.

Unfortunately, most science data systems were developed using differing organizational principles resulting in heterogeneous information architectures. This makes the sear chandre trieval of data products a cross these systems

difficultsinceusersarerequiredtoconnecttoeach individualdatasystemanddealwithdissimilarandoften unfamiliarorganizingprinciples.

Metadataorganizationandmanagementisakeyinformati on architecturalconcept.Metadataengineeringalsoprovides keyorganizationalprinciplesandincludesdatamodeling, dataelementdefinition,anddataresourcedescription.Since metadatahasalsobeenidentifiedaskeyto enabling interoperabilitybetw een heterogeneousdatasystem,the OODTinformationarchitectureincludesareferencemodel fordescribingmetadata.

AkeyfeatureoftheOODTinformationarchitectureisthe useofcommonlyacceptedstandards.WeusetheISO/IEC 11179 [18]specificationforthedefinitionofdataelements inthedataelementregistry.ISO/IEC11179isaframework forthespecificationandstandardizationofdataelements andis widelyaccepted, providing the underlying basis for dataelemen tdefinitionandclassification. The specification defines four data element attribute categories: identification, definitional, representational, and administrative. In addition, we'veadopted use of the Dublin Core [17] elementset,acontentdescriptionmodelthathasreceived wide-spreadacceptanceacrosstheelectronicinformation community. These elements were designed to facilitate the discoveryofelectronicresourcesacrosstheInternetand includecontent, property, andinstantiationresource attributes.

We'vealsoadoptedtheconceptoftheObjectIdentifier (OID)toprovidepersistent,unique,andsimpleobject identification [18].DefinedbytheInternational TelecommunicationsUniona ndadoptedbyInternetSNMP andLDAPcommunities,anOIDispersistentonceassigned anditslifetimeisinfinite.OIDsareuniquebecausethey followadistributedmanagementmodelliketheInternet DomainNameSystem,andaresimplebecausetheyarea seriesofshortintegers.It'salsopossibletotraceanylevel ofanOIDbacktoitsowner.

AnOODTinformationarchitecturalprinciplesuggeststhree phasesinthedevelopmentofmetadataforthepurposeof enablingheterogeneousdatasysteminteroperabi lity. The firstphaseissimplytheidentificationofindividualdata systemdatadictionaries. Datadictionaries describe the elements and associated values used to describe the information or datamodel associated with a datasystem.

Thesecondphaseo fdevelopmentinvolvesthetransforming ofthedatadictionary's native representation into a standard representation based on the ISO/IEC11179 specification. This standard representation is used to describe each data elementidentified in the dictionary. The following XML fragmentillustrates the use of XML to define a space science data element.

<?xml version="1.0"?>
<schema>

```
<dataElement>
  <name>TARGET_NAME</name>
  <identifier>1.3.6.1.4.1.1306.2.10.997
  </identifer>
  <version>2001</version>
  <registrationAuthority>NASA.PDS
  </registrationAuthority>
  <context>Metadata.DataDictionary.Element
  </context>
  <definition>Identifies a target, which maylex bility. The flexibility of XML allowed us to develop a
   a planet, satellite, ring, region, featugenericstructureformanagingresourcedescriptionsfrom
   asteroid or comet. See TARGET TYPE.
  </definition>
  <dataType>CHARACTER</dataType>
  <format>N/A</format>
  <unit>none</unit>
  <maxSize>30</maxSize>
  <obligation>Mandatory</obligation>
  <maxOccurrence>1</maxOccurrence>
 </dataElement>
</schema>
```

Thethirdphaseofmetadatadevelopmentinv olvesthe manualanalysisandcomparisonofdataelementsboth withinandacrossdomainstoidentifyorderiveCommon DataElements(CDEs). This typically involves identifying similar element names or definitions, determining core concepts, possibly gener alizingtheconcept, and determiningkeynamesandaliases. The resultant discipline, application, or enterprises pecific CDEs enable higherlevelsofdatasysteminteroperabilityonceadopted byindividualdatasystems.

TheOODTprofileserverdescribe dinSection metadatadescriptionsofresourcesacrossandwithin distributeddatasystems. The metadatadescriptions are createdfromandvalidatedagainstthemetadataregistryand arepackagedas"profiles."Asm entionedbefore, resources canbealmostanythingincludingsubsystems, services, data volumes, data collections, ordataitems. Examplesofspace sciencedataitemsincludeimages, spectra, and documents.

TheOODTprofileisanXMLdocumentthatuses domain metadatatodescribethedataresourcesthatexistwithina datasystem.Profilespromoteinteroperabilitybetweendata systemsbyprovidingacommonstructureandinterchange languagewithinwhichtodescribesystemresources. The primarypurpose ofanOODTprofileistoprovidea resourcedescriptionthatissufficientforaquerymanagerto determineiftheresourcecanresolveaquery. This subsequently limits the number of resources that will havetoconsiderthequery.Forexample, withins pacesciences,a queryforimagesofJupitershouldnothavetobehandledby theresourcethatmaintainstheMarsGlobalSurveyor images of Mars. A profile can be defined as a proper subset ofthemetadatathatdescribesaresourceandwhichis

sufficienttodeterminewhethertheresourcecouldresolvea query.

InadditiontousingXMLastheinterpeerexchangeformat, we also use XML as the common language for the OODTprofiles. The advantages of XML include (1) superior expressivenesstoHTMLbyallowin ginformation -structure specifications,(2) simplicitycomparedtoSGMLinuseand syntax,(3) wideacceptanceasanElectronicData Interchange(EDI)standard,and(4)themostcompelling,its anysciencedomain.

TheDocumentTypeDefinition(DTD)specificationlisted belowillustratesthebasiccomponentsoftheresource profile.TheDTDspecificationhasthreeparts:theprofile attributes.resourceattributes.andtheprofileelements. We'reusingaDTDspecificationsincethetechnologyis widelysupported.WeareconsideringXML -Schemasince itisnowarecommendationoftheWorldWideWeb Consortium(W3C).

```
<!ELEMENT profiles (preoff)>l
<!ELEMENT profile (profAttributes, resAttributes,
 profElement*)>
<!ELEMENT profAttributes (profId, profVersion?,
 profType, profStatusId, profSecurityType?,
 profParentId?, profChildId*, profRegAuthority?,
 profRevisionNote*, profDataDictId?)>
<!ELEMNT resAttributes (Identifier, Title?,
 Format*, Description?, Creator*, Subject*,
 Publisher*, Contributor*, Date*, Type*, Source*,
 Language*, Relation*, Coverage*, Rights*,
 resContext+, resAggregation?, resClass,
 resLocation*)>
<!ELEMENT profElemententid?, elemName,
 elemDesc, elemType?, elemUnit?, elemEnumFlag*,
 (elemValue* | (elemMinValue, elemMaxValue)),
 elemSynonym*, elemObligation?,
 elemMaxOccurrence?, elemComment?)>
```

TheprofAttributessectionoftheXMLdocumentcontains theattributesof theprofileitself. The profIdattribute providesasystem -wideuniqueidentifierforaprofile instanceandiscurrentlybeingimplementedasanOID. The prof Title and prof Descattributes provide descriptions of theprofilewhereprofTitleismoreterse andappropriatefor morefrequentdisplayinuserinterfaces. The prof Type attribute, defined further below, identifies the profile type. TheprofDataDictIdattributeprovidestheidentifierofthe controllingdomaindatadictionary. The prof Child Id attributes allow for the creation of a profile hierarchy.

TheresAttributessectionoftheXMLdocumentgenerically describestheresourceusingthe15recommendedelements of the Dublin Core initiative. A description of this initiative andtheelementsisa vailableathttp://purl.org/DC/.We've addedthreeadditionalresource attributes.TheresContext elementidentifiestheapplicationenvironmentordiscipline withinwhichtheresourceoriginatesandisderived taxonomyofsciencedisciplines.Fore xample, NASA.PDS.Geoscienceisusedtoindicatethattheresource is associated with the Geoscience node of the Planetary Data System. Theres Location provides the location of the resource, typically represented as a URL. Finally, the resClasselementi dentifiestheresourcewithinataxonomy ofresourcetypes.Examplesvaluesare system.productServer,application.interface,data.granule anddata.dataSet .

The profile elements ection of the XML document describesthedatacontentthattheresourcemanages usingdomain specificmetadata.Forexample,thePlanetaryDataSystem (PDS)maintainsaninventoryofallsciencedatasetsthat havebeenarchivedinthesystem. Withinthis inventory, the datasetsareindexedoninstrumentandtargetbody identifiers involved in the observation. These elements wouldbeincludedinaprofile's element section as illustratedbelowandindicatethatthisdatasetcanresolve queries that have PDS in strument and target identifiers asqueryconstraints.

Asillustratedin theXMLfragmentbelow,eachdata elementisdefinedusingthemeta -attributeselemId, elemName.andelemValue.Thesemeta -attributesare consistentwiththoseintheOODTdataelementregistryto allowforreferenceandvalidation.

```
file>
 fAttbutes>
  fId>1.2.3.4...
  fStatusId>ACTIVE/profStatusId>
</profAttributes>
<resAttributes>
  <Identifier>1.2.3.5...</Identifer>
  <Title>Viking Orbiter 1...</Title>
  <Format>text/html</Format>
  <Languagen</Language>
  <resContext>NASA.PDS</resContext>
  <resClass>data.dataset</resClass>
  <resLocation>http://pdscat...</resLocationpexistencesincethelate1980s,thePDSdeveloped</pre>
 </resAttributes>
 profElement>
  <elemId>1.2.3.6...
  <elemName>INSTRUMENT N...
  <elemType>ENUMERONI
  <elemValue>VISUAL IMAG...
</profElement>
 fElement>
  <elemId>1.2.3.7...
  <elemName>TARGET_NAME
  <elemType>ENUMERATION</elemType>
  <elemValue>MARS</elemValue>
  <elemValue>PHOBOS</elemValue>
  <elemSmyonym>ADS.OBJECT ID</elemSynonym>
```

```
</profElement>
</profile>
```

Thefinalphaseofprofiledevelopmentinvolves implementinginstancesofprofilesforspecificdomain resources. Asisevident, the success of the distributed resourcelocationconceptisdepen dentontheexistenceof domainmetadatacapturedinrepositoriessuchasdata (element)dictionaries. Within such privileged domains, the registrationofresourceswiththeserviceisreadily accomplishedbyextractingthenecessarymetadatafromthe domain's metadata repository, creating the resource profile, andthenregisteringtheprofilewithaprofileserver. This enablesthesuccessfullocationofresourceswithina domain.

Supportinginteroperabilitybetweenresourcesfrom differentdomainsisst ronglydependentonmetadata compatibility, or how well the metadatas pans the domains. Forexample,tworelateddomainssuchasplanetaryscience andastrophysicsbothassociateoneormoretargetbodies withmostdataproducts. However, unless the same identifiersareusedforaspecifictargetoramapping betweenidentifiersisdetermined, the attribute will not supportresourcelocationmuchlessinteroperabilityacross thedomains.Infact,asmoresophisticated interoperability—suchasdatatransfo rmationand correlation—arerequested, deeperlevels of metadata compatibilitywillberequired.Forexample,onceatarget bodyisidentified, sufficient metadata must be available for coordinatesystemconversion.

Thedevelopmentofcross -disciplinemet adataistypicallyan arduous manual effort performed by teams of domainexperts. Howeverthere are research efforts focusing on automaticmethodsforidentifyingcommonmetadataacross domains.

5. APPLICATIONTOTHE PLANETARY DATA **SYSTEM**

ThePlanetaryData System(PDS)managesandarchives planetarysciencedataforNASA'sOfficeofSpaceScience. -early on—astandardsarchitecturethatincludedaformal enterprisemodel, ameans for collecting and associa ting metadatawithsciencedataproducts, and apeerreview processforensuringdataandmetadatavalidity. This active sciencedataarchivecurrentlyhasoverfiveterabytesofdata stored and distributed on CD and DVD media. Thestandardsarchitecture hasproventobethesinglemost importantelementinmaintainingconsistencyacrossthe varioussciencedomainsrepresentedintheplanetaryscience communityandthesixgeographicallydistributedscience disciplinenodesofthePDS.

Theprimaryfocuso fthePDSto -datehasbeenthecreation ofalong -termsciencedataarchiveonphysicalmedia.PDS providescommunityaccesstothedataproductsprimarily

through the appropriate distributed science discipline node. Forexample,thePDSimagingnodeisr esponsiblefor imagingdataproductsandbuildssearchandretrieval interfaces for specific datasets as requested by their users. Similarly, other discipline nodes are responsible for the data withintheirdomainandbuildsearchandretrievalinterface S asneeded. The guiding philosophy was that scientists workingintheplanetarysciencecommunitywouldknow whatnodeorinterfacecouldprovidethedataofinterest. Thisdevelopmentapproachresultedinheterogeneous searchandretrievalsystemsanda sthenumberofdatasets increased, an exponential increase in the number of point -topointinterfaces.

TheadventoftheWeb, ahuge increase in potential users, and the larged at a volumes proposed for future solar system explorations missions has re-focused attention on distributed data management is sues. These include an eed for an integrated view of all the data in the archive, an integrated spatial/temporal access a cross the entire archive, and infrastructure support for science analysis and knowledged discovery.

thePDSdevelopedtheprototypeDistributed In1998. InventorySystem(DIS)toinventorytheresourcesacross the PDS. This prototype took advantage of the wealth of metadatainthearchivetodescribedbothdataandnon -data resources. The DIS described each resource in a labelusing pertinentmetadatafromthearchive.TheDISusest he ObjectDescriptionLanguage(ODL),akeyword/value language,toencodethemetadata.URL/FTPlinkswere includedforallonlineresources. The labels wer ecollected intoasimpledatabase.DISengineersdevelopedasimple PerlCGI -scriptsearchenginetosearchthelabeldatabase. ResultsareprovidedasanHTMLdocumentfordisplayina browser.

ThisprototypewasasuccessandprovidedthePDSwith severallessons.First,iti dentifiedthePDSmetadataasthe keytotheintegrationproblemsinceitwassufficientto identifyanddescribebothdataandnon -dataresources, providedthenecessarysearchattributesforfinding resources, and was sufficient todeterminewhetheraquery couldberesolved by the resource. In addition the prototype showedthattheincreaseofpoint -to-pointinterfacescould bereducedtoalinearfunction. The prototype also made certainlimitationsobvious. Thesewere that O DLwasnota universalinterchangelanguage, resourceheterogeneitywas stillvisibletotheuser,navigationandaccesswaslimitedto HTTPlinks, and noded at a system interoperability is not supported.

Inlate 1999, PDS and OODT staffmetto discuss the applicability of OODT's dataman agement technology to PDS is sue sand PDS's information architectural concepts (datamodeling and metadata standards) to OODT goals.

One result of this collaboration is called the Next Generation DIS. The goals of this news ystemin cludes adopting a common interchange language (XML), hiding

heterogeneityusingencapsulation,usingmessagepassing protocolinacomponentbasedsystemarchitecture, maintainingsystemlocationindependence,andprovidinga scaleableandextensi blesolution. Figure 7illustratesthe proposedsystemwhichwillbeprototypedusingtheMars GlobalSurveyor(MGS)sciencedataarchives.

Next Generation DIS

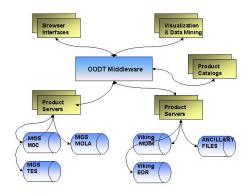


Figure 5 -NextGenerationDIS —MarsGlobalSurveyor Prototype

OODTprofileswillreplacetheODLresourcelabelsofthe initialDISsystem.AportionofanOODTresourceprofile fortheVikingMarsImagingDigitalImageModeldataset (MarsMosaiccreatedusingoriginalVikingimages)is illustratedbelow.Thisprofileprovidessufficient informationforaqueryhandlertodeterminewhetherthe datasetcouldresolveanincomingquery.Inaddition,this profilecouldalsobeeasilymodifiedtodescribeaproduct serverthatservesindivid ualimagesbysimplychangingthe resourcedescriptionsincethemetadataencodedinthe profileelementsectionisapplicabletobothresources

```
cprofile>
 fAttributes>
  fId>1.3.6.1.4.1.1306.2.102
  </profAttibutes>
 <resAttributes>
  </Identifer>
  <Title>VO1/VO2_MARS_VISUAL</Title>
  <Format>text/html</Format>
  <resContext>NASA.PDS</resContext>
 <resClass>data.dataSet</resClass>
  <resLocation>http://pdsproto.pplgonas..
  </resLocation>
 </resAttributes>
 profElement>
 <elemName>TARGET NAME</elemName>
  <elemType>ENUMERATION</elemType>
  <elemValue>MARS
 </profElement>
 profElement>
 <elemName>MAXIMUM_LATITUDE</elemName>
 <elemType>REAL</elemType
```

<elemMinValue>.50000</elemMinValue>
 <elemMaxValue>90.00000</elemMaxValue>
 </profElement>
</profile>

Theinitialgoaloftheprototypewillbetosupportthe creationofadatacoverageplotasillustratedin Figure 6.I n thisplot,indicatorsshowingthegeographicallocationof MGSimaging(red),altimeter(blue),andthermalemission spectrometer(green)dataproductsareoverlaidontoaMars imagemosaicfromtheVikingMarsimagedataset.Four datasetsatseveral geographicallydistributeddatasystems willbeaccessedtoproducetheplot.

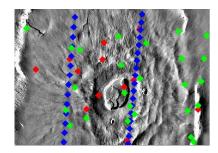


Figure 6 -MarsGlobalSurveyorCoveragePlot

6. APPLICATIONTOTHE DEEP SPACE NETWORK

Establishedin 1958, NASA's Deep Space Ne twork (DSN) is a global network of an tennast hat communicate with spacecraft. The DSN facilitates interplanetary spacecraft missions and several Earth - orbiting missions. It is currently made up of three facilities across the world: at Goldstone, in California's Mojave Desert; near Madrid, Spain; and near Canberra. Australia ³.

TheDSNplaysanintegralroleinthesuccessofmissions. Amongstafewofitsdutiesare:acquiringtelemetrydata fromspacecraft,transmittingcommandstospacecraft, gathering sciencedata,andtrackingspacecraftpositionand velocity³. However, these tasks are hindered by the distributed and heterogeneous nature of the DSN. These obstacles complicated at a sharing and mining, data modeling, an dthecreation of a common information architecture.

TheOODTframeworkdirectlyaddressesthekeyissuesof theDSN.OODT'semphasisonimplementationofbotha componentframeworkalongwiththeidentificationoflocal datadictionariesandmodelsfor each system provides an excellent vehicle forestablishing an information architecture for the DSN. In addition, OODT's design principle of hidinglocal system interfaces through data abstractional lows for the implementation of a common information architecture. OODT allows for data elements

tiedtolocalinformationsystemstobetranslatedinto commondataelementsdescribedfortheentireDSN.Its extensiblearchitectureprovidestheDSNwithwhatcould beconsidered "plugandplay" datamanagement. In addition, this architecture provides ame ans by which data and local agents can be located, retrieved or executed on any number of systems and machines using a common XML exchanges tructure as defined in section 3. This distoributed computing model fits very nicely into the overall DSN architecture.

Asproofofconcept, we created and launched a prototype fortheDSN.Thefirstphaseofthisprototypeinvolved creating the ``Mission and Ground Asset Database." Thisdatabasepr ovidesabasicdescriptionofmissions, ground stations, spacecraft, and ground antennae. It was created to provideanonlinereferenceofkeymissionparametersand valuesandtoprovidealongtermvisionfortyingtogether capabilities of ground statio nsaround the world. Figure 7 showsahighlevelentity/relationshipdiagramofthe database. This model will allowengineers to discover possibleconstraintsofthecurrentgroundsystemsandor missions.Inaddition,thedatab asecouldbeusedto understandhowtooptimizeworkflowsinandaround groundstations.

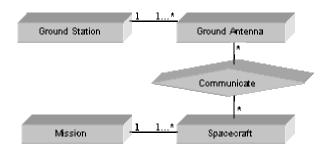


Figure 7 -ERDiagramofMissionandGroundAsset
Database

Theimplementation of the database in conjunction with the OODTarchitectureal lowedforthegroundstation informationcapturedtobepublishedandexchangedover theInternetusinganXMLinterface.Notonlyisthe databasenowsearchablewithinthearchitecture,itcanbe discovered by other systems by passing the XML structuretotheOODTsystem.OODTusedaproductservertoquery and return the results of the database using the OODTXMLquerydefinition. The products erver was defined in a profilewhichdescribesthemissionsandgroundstations thataremaintainedbythedat abase. This allows for a query ofmissioninformationtofirstlocateadatabasewhich containsthedesiredinformation.Lookingforward,thiswill allow formultiple databases to be tied together fromdisparategroundassetdatabasesaroundtheworldif standardXML structure can be defined for querying thatdata. An excerpt of the XML structure from the profile canbeseenbelow.

a

³http://deepspace.jpl .nasa.gov/dsn/

```
file>
    fAttributes>
        fIdb. 3.6.1.4.1.1306.5/profId>
          fVersion>1.0fVerision
        </profAttributes>
    <resAttributes>
   <Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PRODUCT_SERVER</Identifier>JPL.TMOD.PROD
        <Title>Mission and Ground Asset DB</Title describethefunctionstheycancarryout.Profilepeers
        <Format>text/html</Format>
        <resContext>NASA.JPL.TMOD</resContext>
        <resClass>application.interface</resClass}retrievalservices.Inthismodel,auserwould"</pre>
        <resLocation>http://velcro.jpl.nasa.gov..
        </resLocation>
   </resAttributes>
    fElement>
        <elemName>MISSION NAME
        <elemType>ENUMERATION</elemType>
        <elemValue>Cassini</elemValue>
        <elemValue>Galileo</elemValue>
   </profElement>
   profElment>
        <elemName>GROUND_STATION_LOCATION/elemNa
        <elemType>ENUMERATION</elemType>
        <elemValue>Goldstone</elemValue>
        <elemValue>Canberra</elemValue>
    </profElement>
</profile>
```

JPL's Interplanetary Network and Information SystemsStandardsProgramh asinitiatedseveraleffortstoexplore the use of XML as a mean stoest ablish common datadefinitionandaccessmechanismsforareferencemission informationarchitecture. This architecture is being worked withsupportfromtheConsultativeCommitteeon Space DataSystemsthroughanXMLWorkingGroup,andthe ObjectManagementGroup'sSpaceDomainTaskForce. Thegoalistoestablishworld -widemechanismsfor exchangingmissioninformationaswellassupportingthe useofXMLthroughoutthemissionlife cycle.OODTis continuingtoworkonsupportingthisvisionandisnow workingtoallowotherdatabasestobequeriedusing OODT'sXMLcomponents.XMLandOODTplayacritical roleinhelpingboththeDSNandothereffortsfromNASA andotheragenciesesta blisharchitecturesforthecapture, location, and exchange of critical mission information acrosstheboundariesofdistributedandheterogeneous systems.

7. Future

We'verealized the physical deployment of the current OODTframeworkusingCORBAasthenet worksubstrate. CORBA's object model represents servers as interfaces with methods, while clients call those methods. The logical deployment, however, is different. It is a set of interconnectedpeersthatpassaroundaqueryobject.Each peerexamines thequeryandmayaddanyresultsthatare relevant. The query service, described in Section simulatesthisP2Parchitecturebycrawlingthenetworkof

serversonbehalfoftheuserandmaintainingthequery object.

Byrep lacingtheCORBAsubstratewithaP2Psubstrate,we canbetterleverageP2Pfeatures.Forexample,thequery servicecompletelydisappearsbecauseitsroleisredundant inaP2Parchitecture.InP2P, therearenoclientsorservers, blishserviceadvertisementsthat wouldpublishadvertisementsforresourcelocationservices. Productpeerswouldpublishadvertisementsfordata queryintoasystem.Serviceadvertisementswouldindicate whatpeers would be capable of handling the query. Results wouldcomebacktotheuser'speerasynchronously; and the usercanterminatethequeryassoonass/hereceivesthe desiredr esults.

Furthermore, aP2P substrate addresses scalability. In the CORBAmodel, the query service must wait its turn when runningaqueryonbehalfofauseratanoverburdened profileorproductserver.Inthepeermodel,apopulardata ne > centercanbele veragedbyprovidingmorethanonepeerat thecentertohandlerequests. These peers may be located onmultipleCPUsandnetworkinterfacestomaximize concurrencyandavailability. In essence, a system administratorcanspeedupasitebystartingmore peers.

Inaddition, wearefocusing on developing new components inthearchitecturetomanagemetadatabycreatinga metadataservice. The purpose of the metadataservice is to provideaseriesofregistriesthatallowforthemanagement ofdatadiction aries,dataelementsandresourcedescriptions. Thisservicewillbedesignedaroundthreemajor capabilities:capturingdatadictionaryschemas,capturing dataelementsandcapturingresourcedefinitionsfordata systems.datasets.anddataproducts.We tothevalidationofprofilesfordescribingsystemresources.

OODTisalsocontinuingtoworkacrossdisciplines.We are working with the National Institutes of Health to explorethe use of OODT for capturing and sharing researchproducedindisparateresearchlaboratoriesandareworking with JPL Institutional Computing and Information Servicesprogram(ICIS)toexploreinfusionofdatamanagement services and information architectures for describing, buildingandoperatingdistribu teddatasystemsatboththe institutionalandprojectlevel.

ThesuccessofOODThasinmanywaysbeendrivenbythe desiretoshareinformation. Whiletherearemany technical, socialandlegalissuesrelatedtodatasharingand interoperability,the prospectofderivingnewknowledge throughdatacorrelationisimmenselyhigh.Inaddition, providingstandarddatamanagementframeworksfor buildinginformationmanagementsystemsthatcan dynamicallyplugintolargerinformationmanagement systemsisc ertainlyagrandvision.Onethingisclear: scienceandengineeringplantocontinuetoproducelarge

3,

volumesofdatathatareordersofmagnitudelargerthan thosealreadycaptured. Theneedtobettercapture, organize, describe, reuseandshareinfor mationwithin and acrossdisciplines will continue to be keyinformation management needs and research topics in the foresee able future.

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